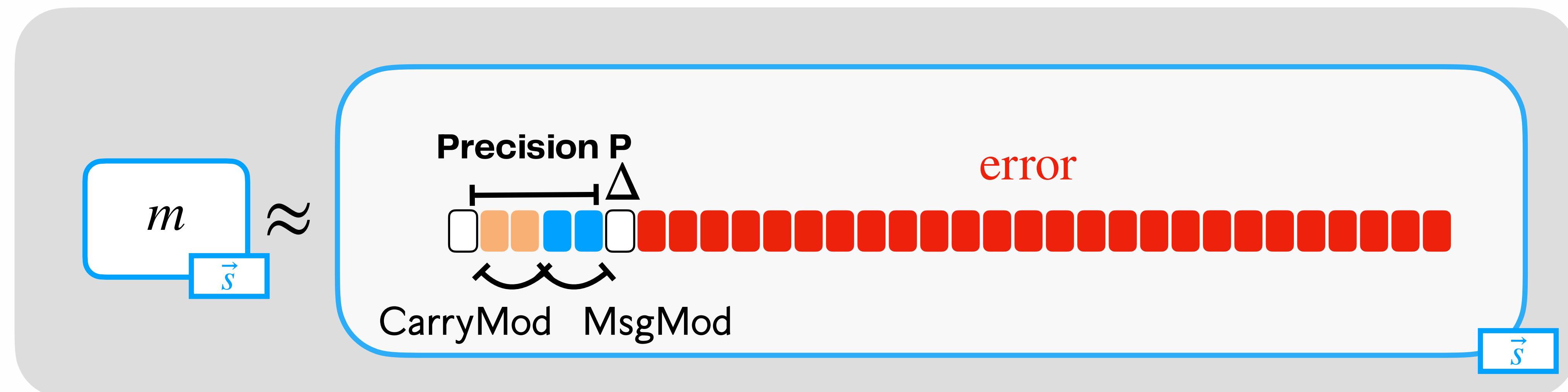
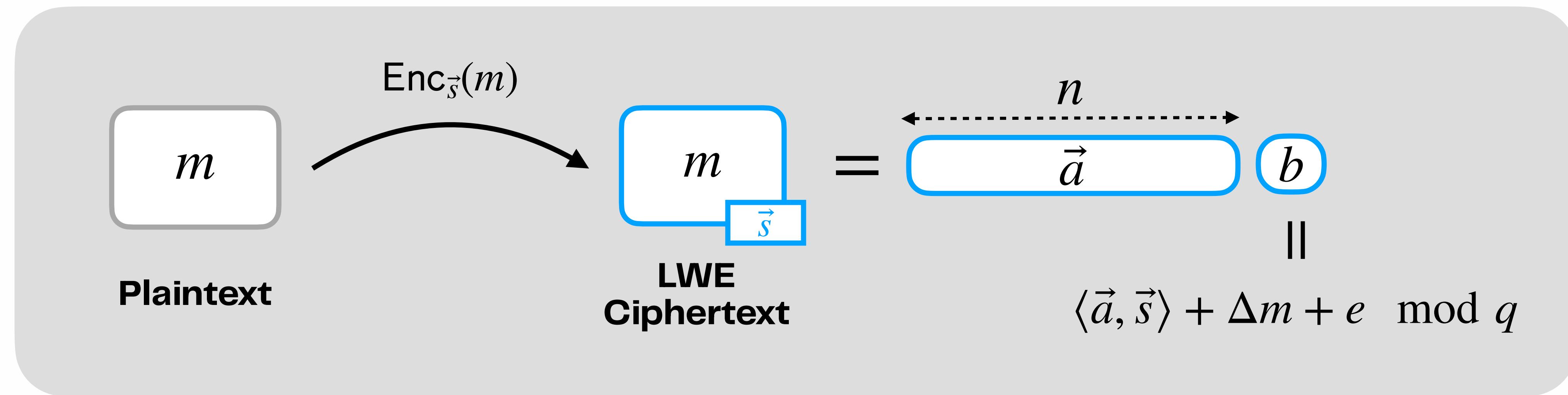


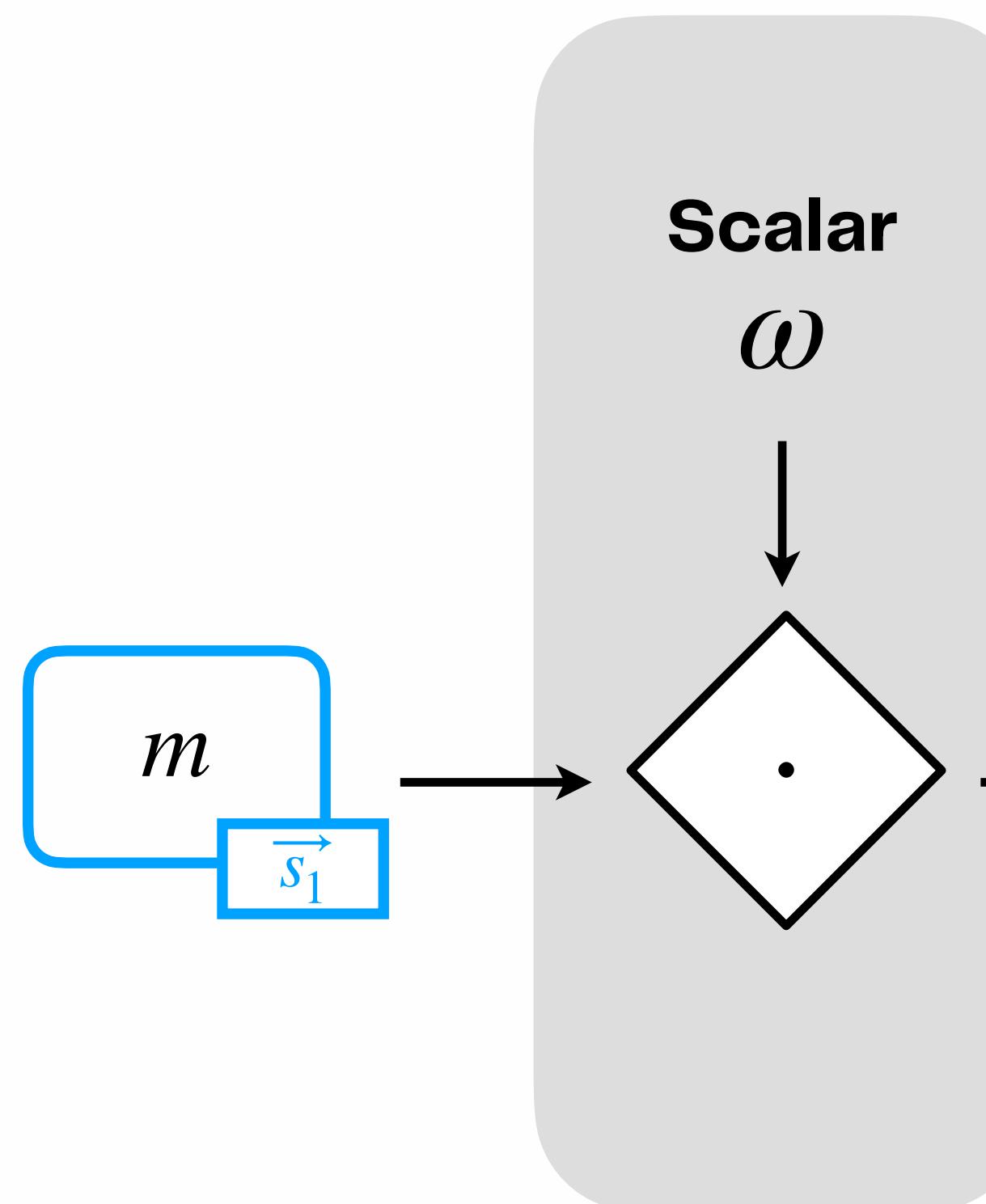
TFHE Simplified: A practical Guide to Integer Arithmetic and Reliability

TFHE Ciphertexts

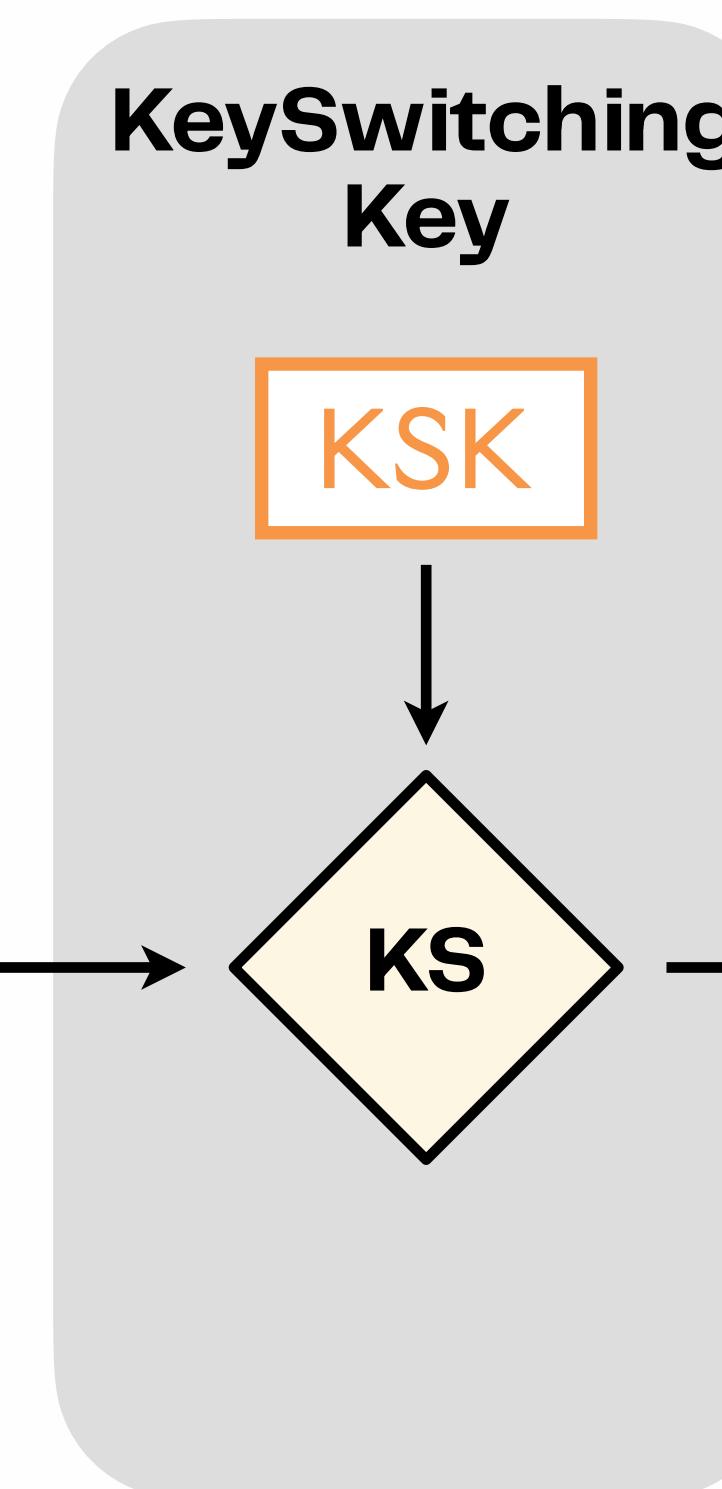


TFHE Toolbox 1/2

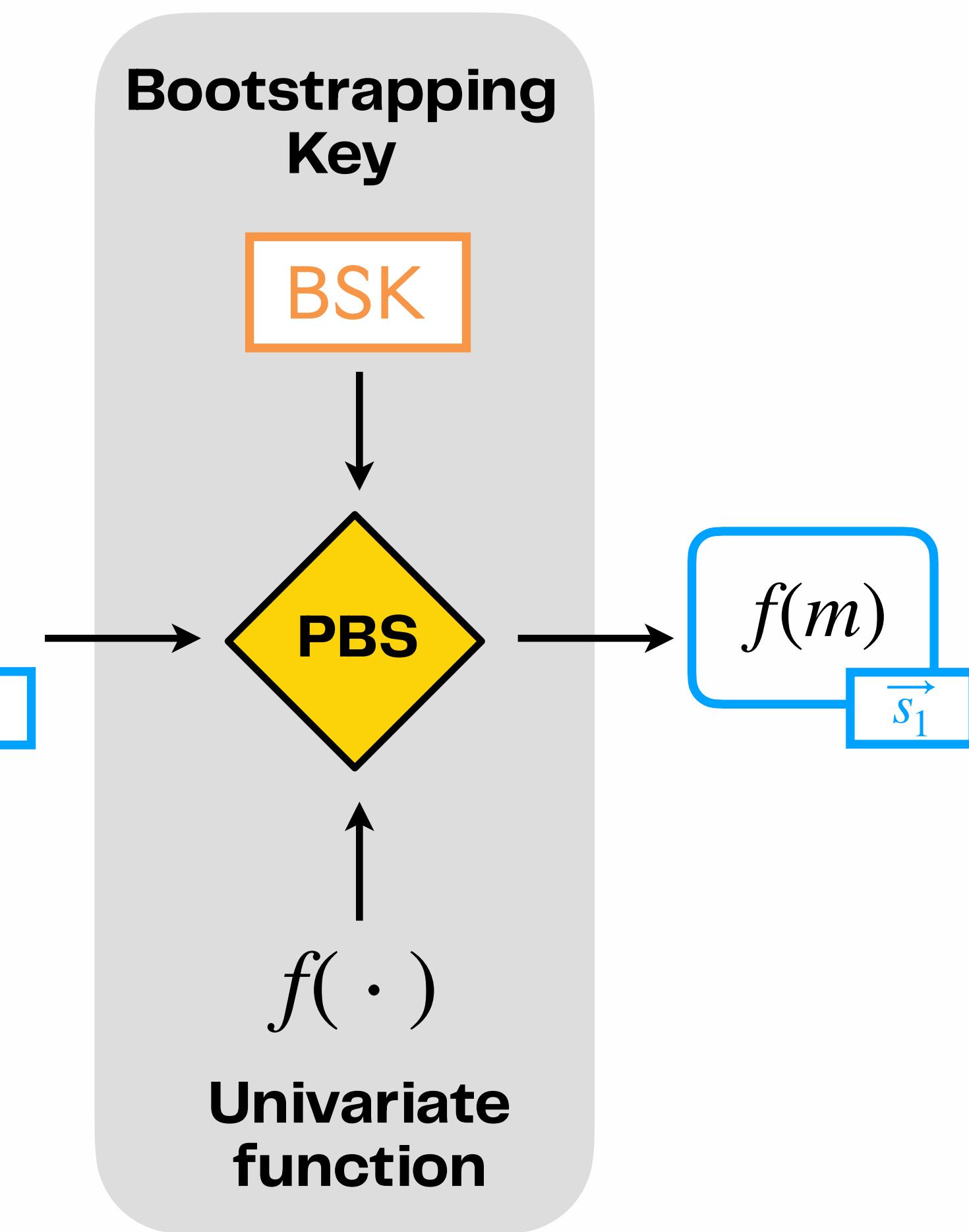
Linear Operations



KeySwitch (KS)



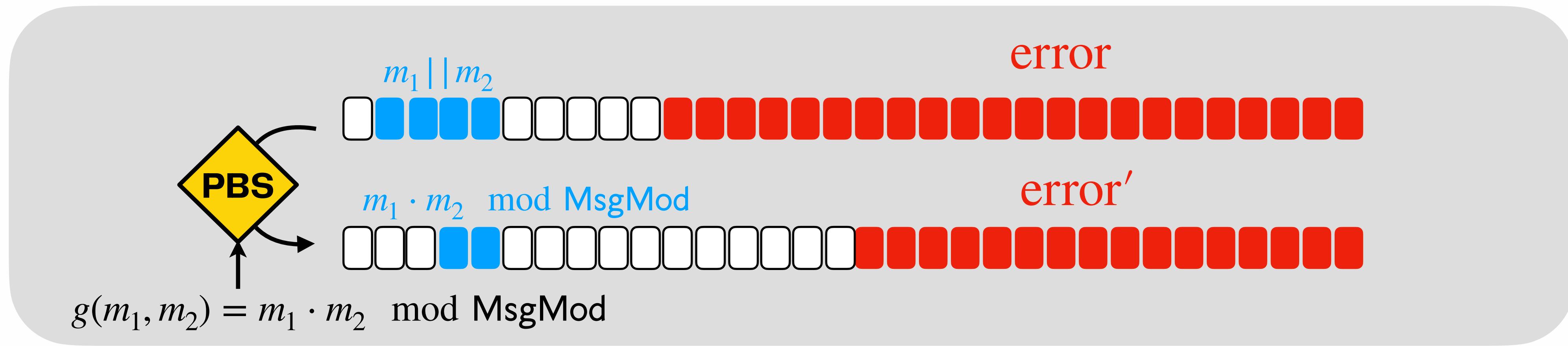
Programmable Bootstrapping (PBS)



Small messages only:
 $|m| \leq 10$ bits

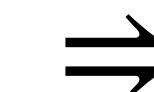
TFHE Toolbox 2/2

Bivariate PBS: Multiplication Example



Splitting cleartext space in two parts:

$$f(m) = g(m_1, m_2)$$



Constraint:

$$\text{MsgMod} \leq \text{CarryMod}$$

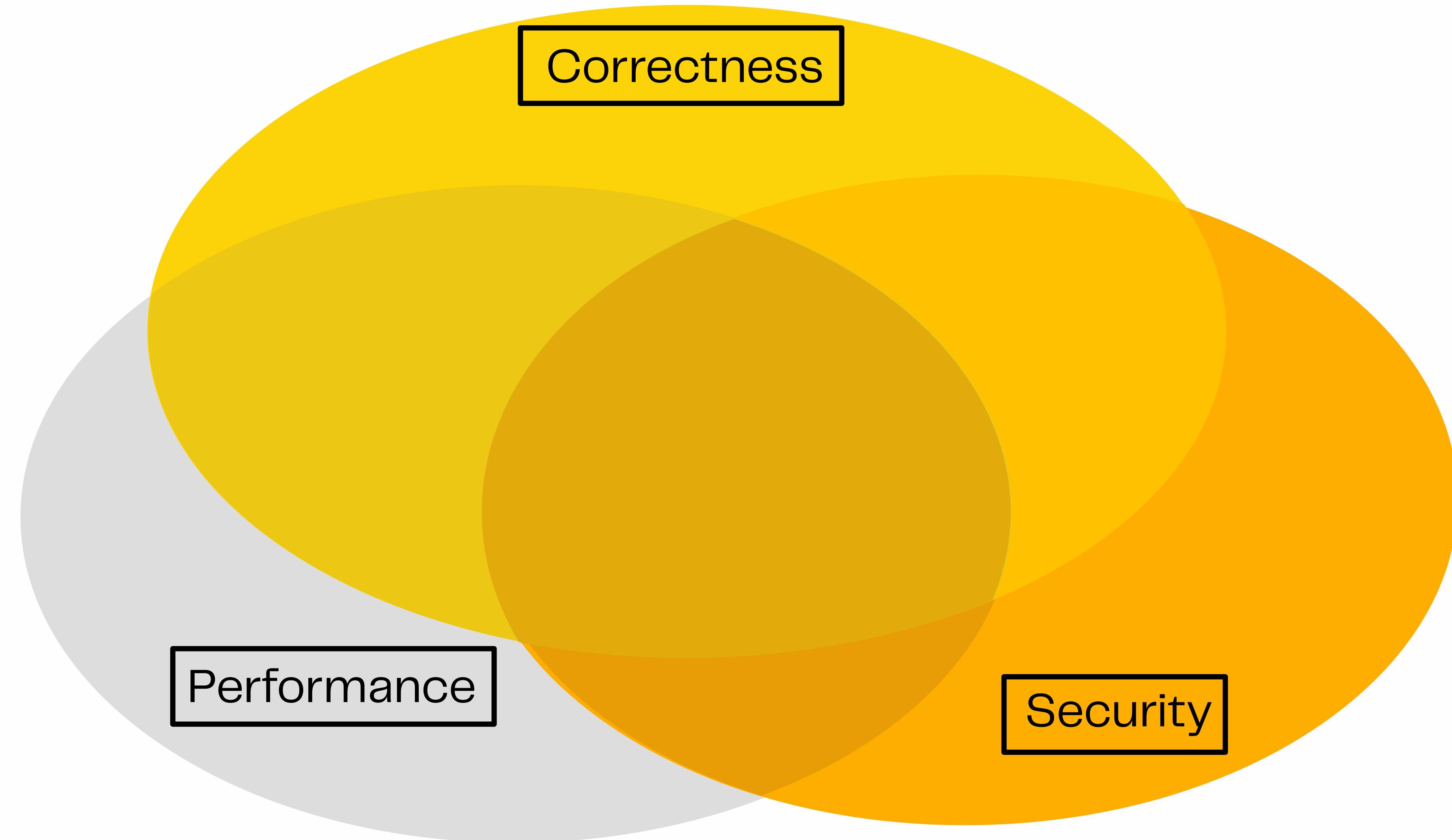
A few parameters...

Correctness

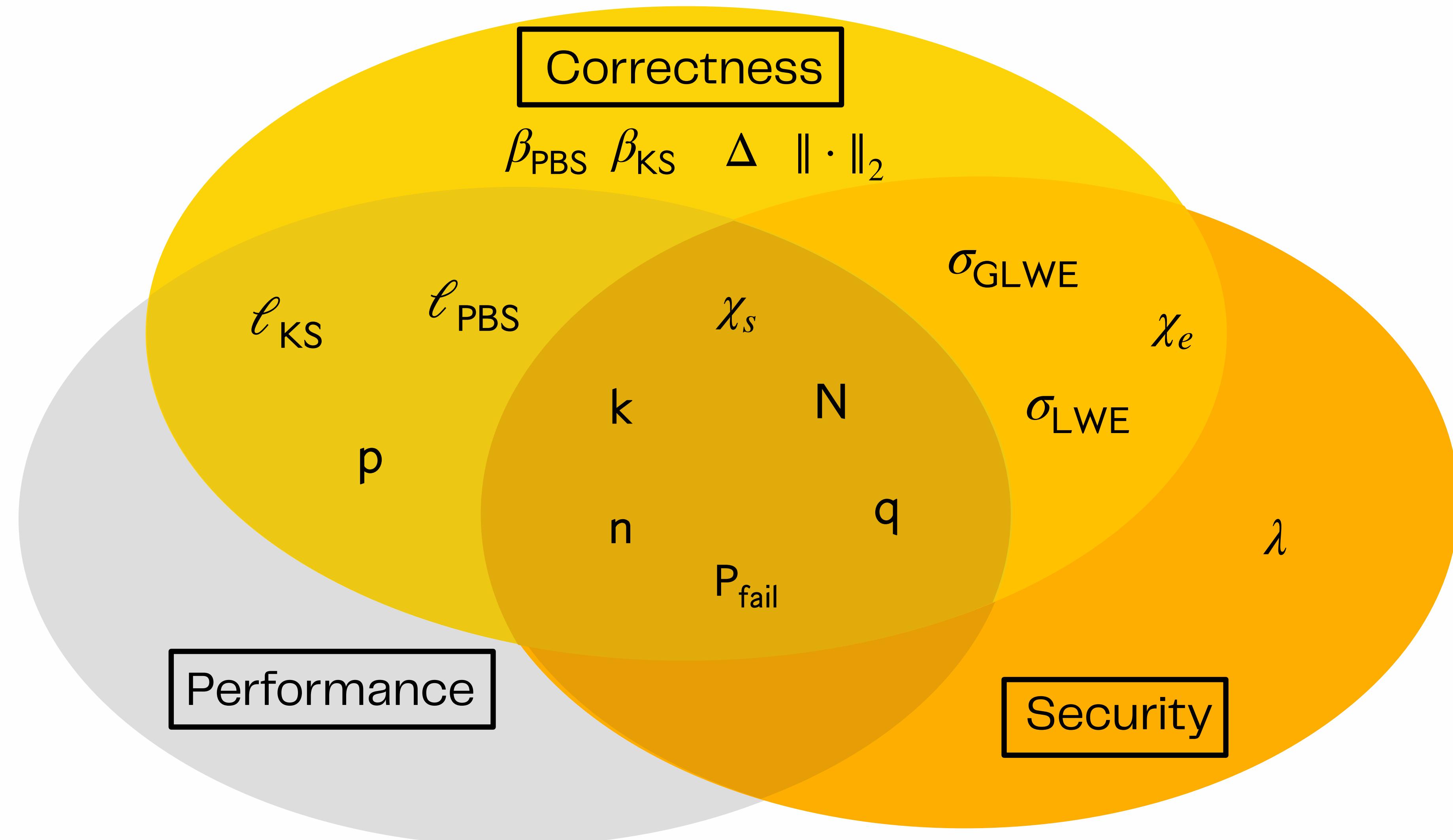
Performance

Security

A few parameters...



A few parameters...



Parameter Oracle

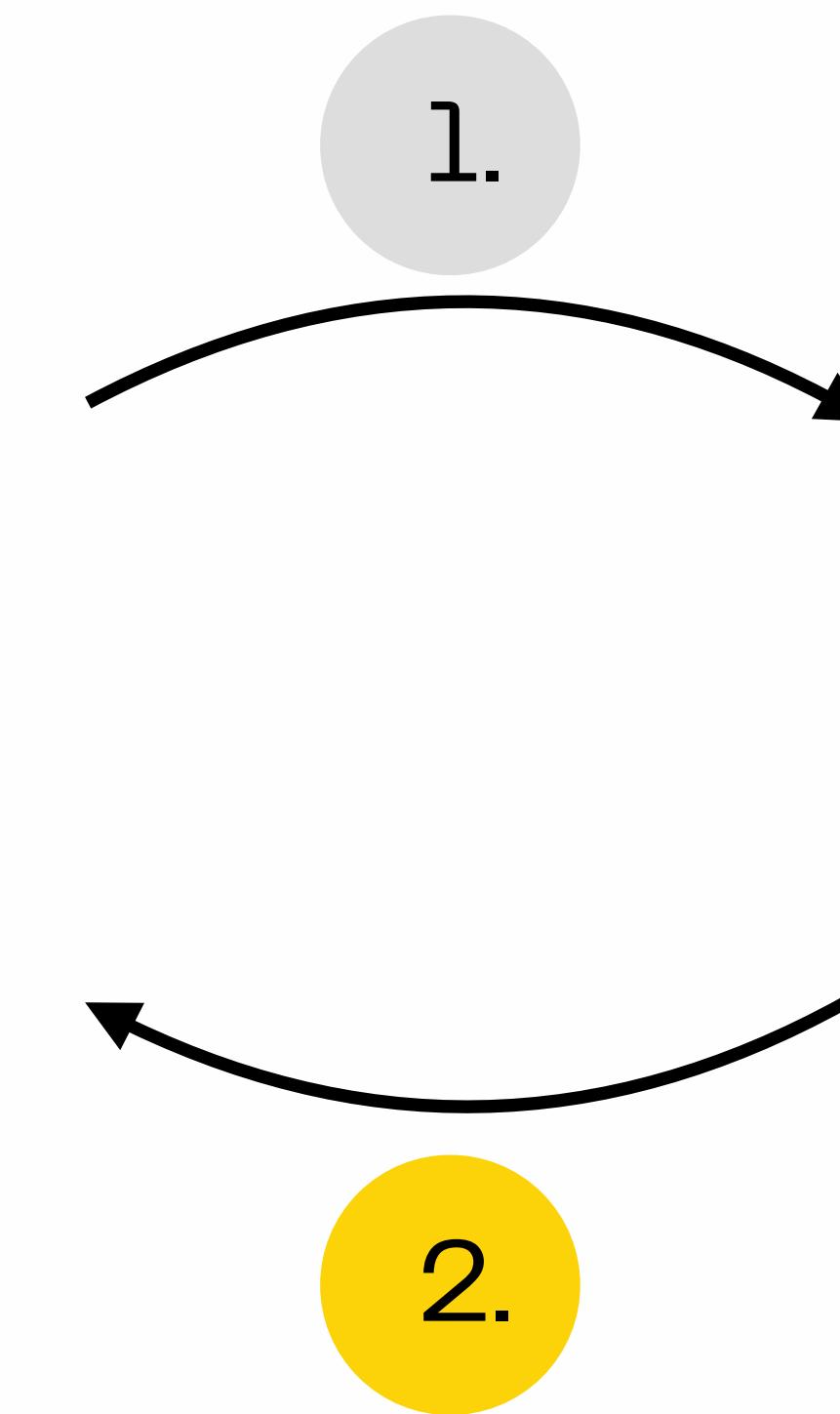
Client

Message precision P

Linear operation number $\|\cdot\|_2$

Error Probability P_{fail}

Security Level λ



Oracle*

β_{KS}	ℓ_{KS}	β_{PBS}	ℓ_{PBS}
n	k	N	
σ_{LWE}	σ_{GLWE}		
χ_s	χ_e		
Δ	q		

FHE constraints

Security, Correctness
and Performance should
be **assured**

Cryptographic parameters
must be **correctly** chosen for
the users

By default, precision of
messages \leq 10 bits

Extension needed to support
u16, u32, u64, ...

How can we abstract **FHE programming**
from its cryptographic complexity to match
the simplicity of **traditional coding**?

Summary



A Foolproof Parameter Oracle



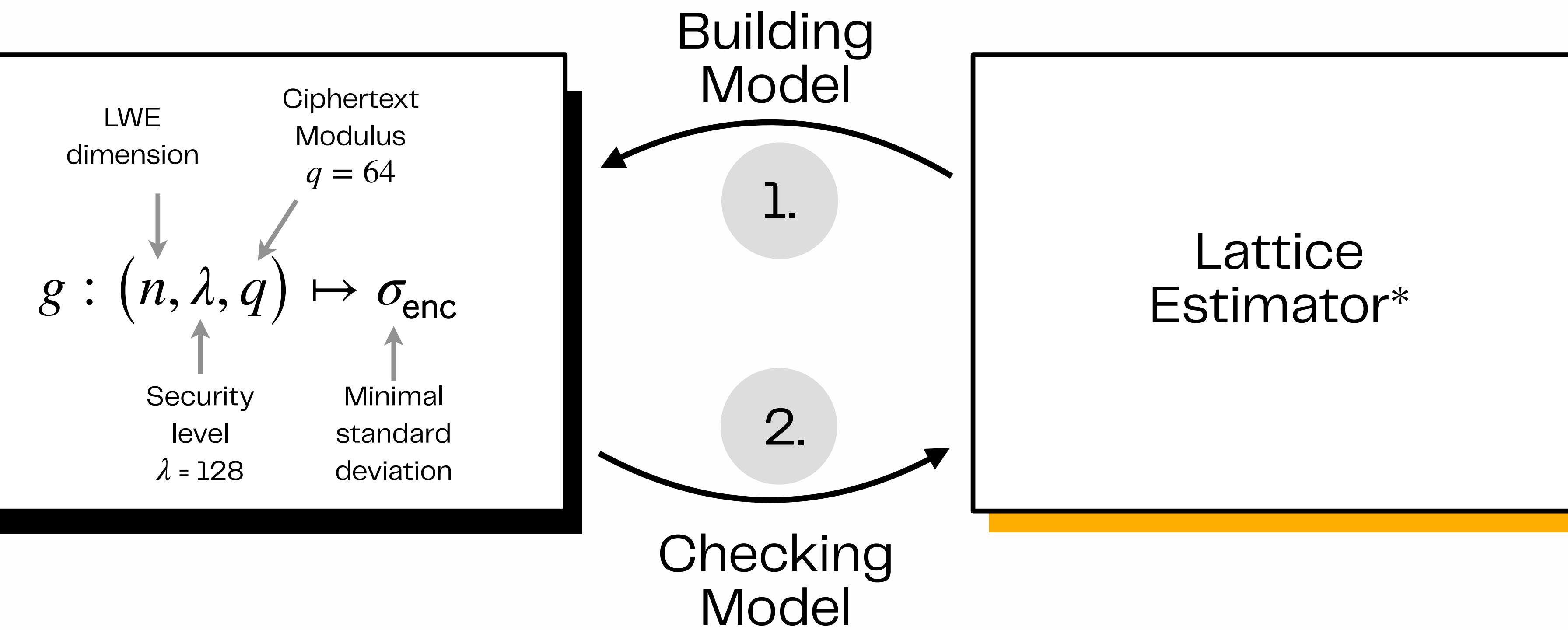
Building Homomorphic Integers



Benchmarking FHE

A Foolproof Parameter Oracle

FHE Requirements

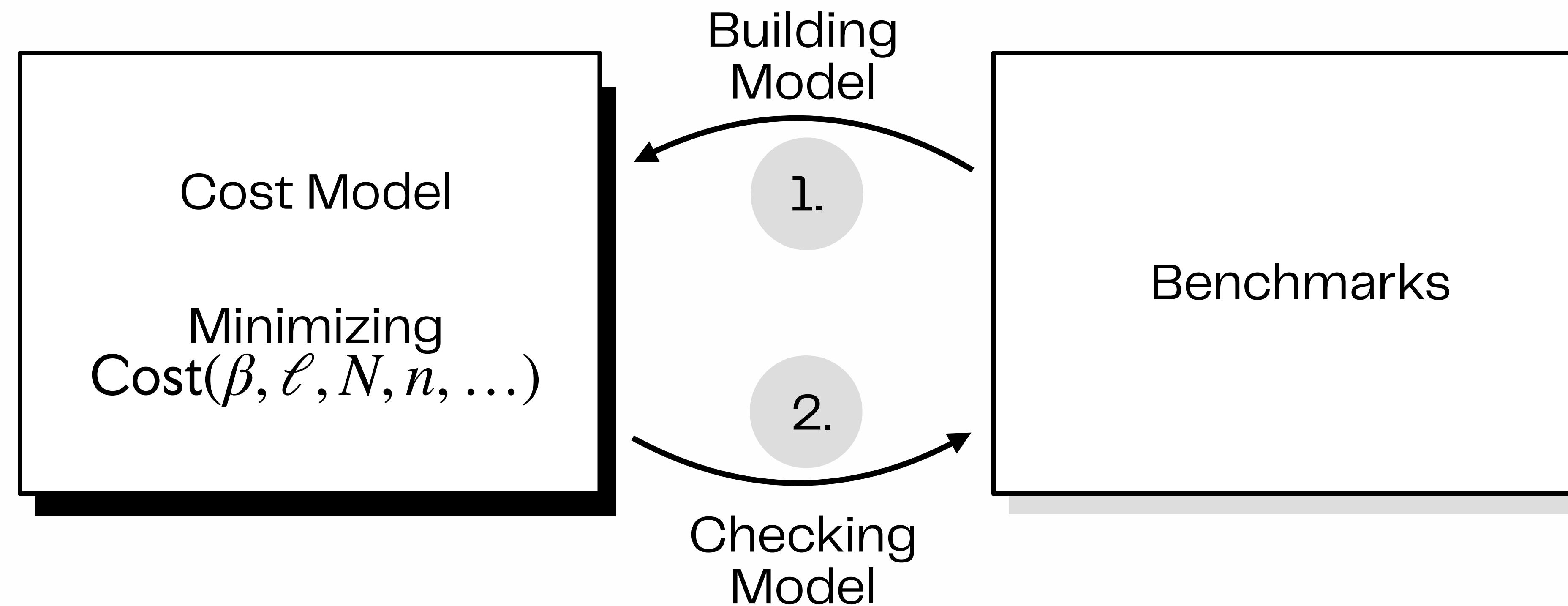


FHE Requirements

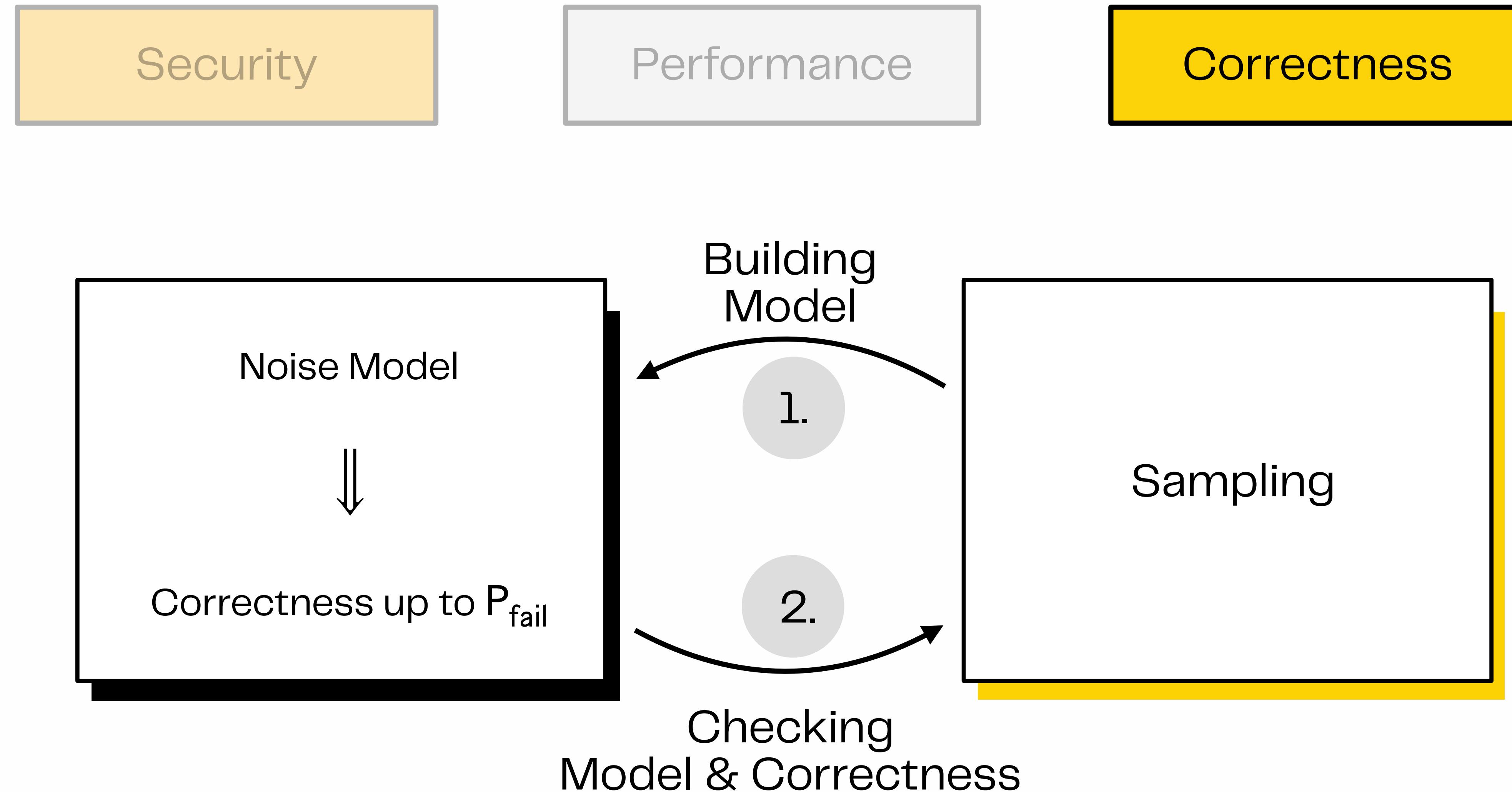
Security

Performance

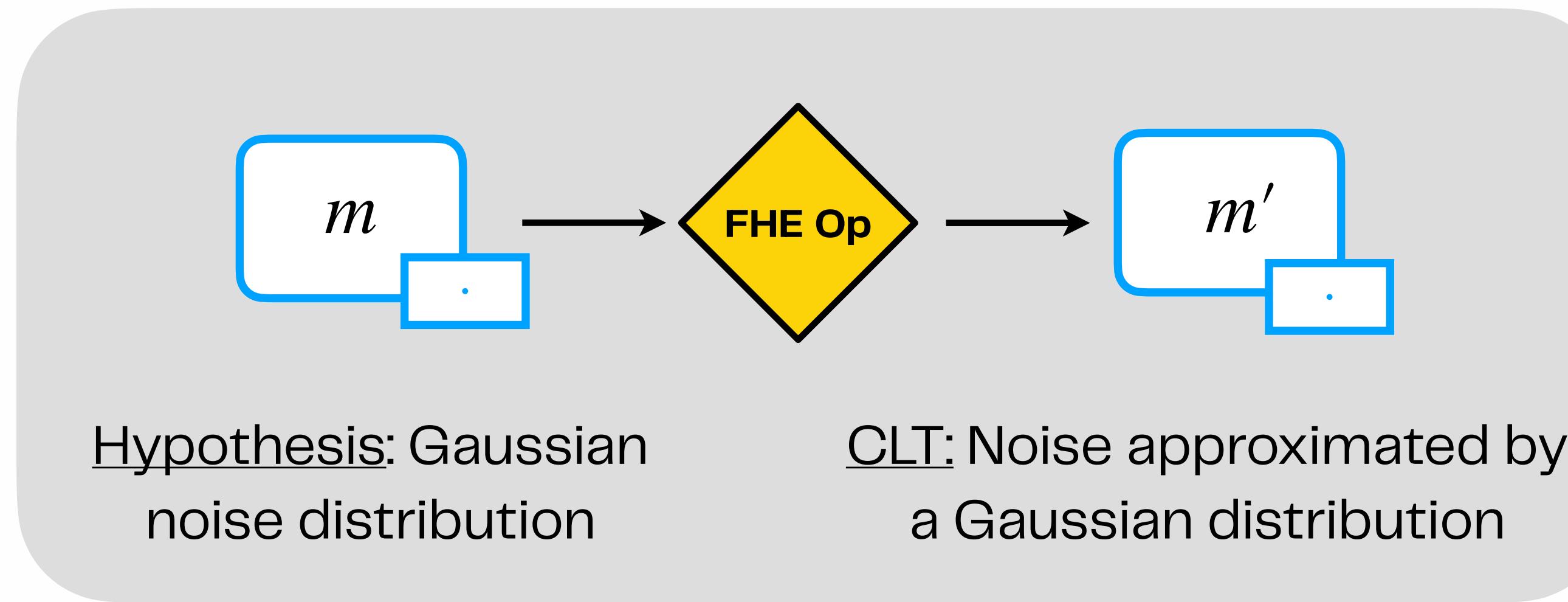
Correctness



FHE Requirements



Checking Noise Model



1.

Normality check

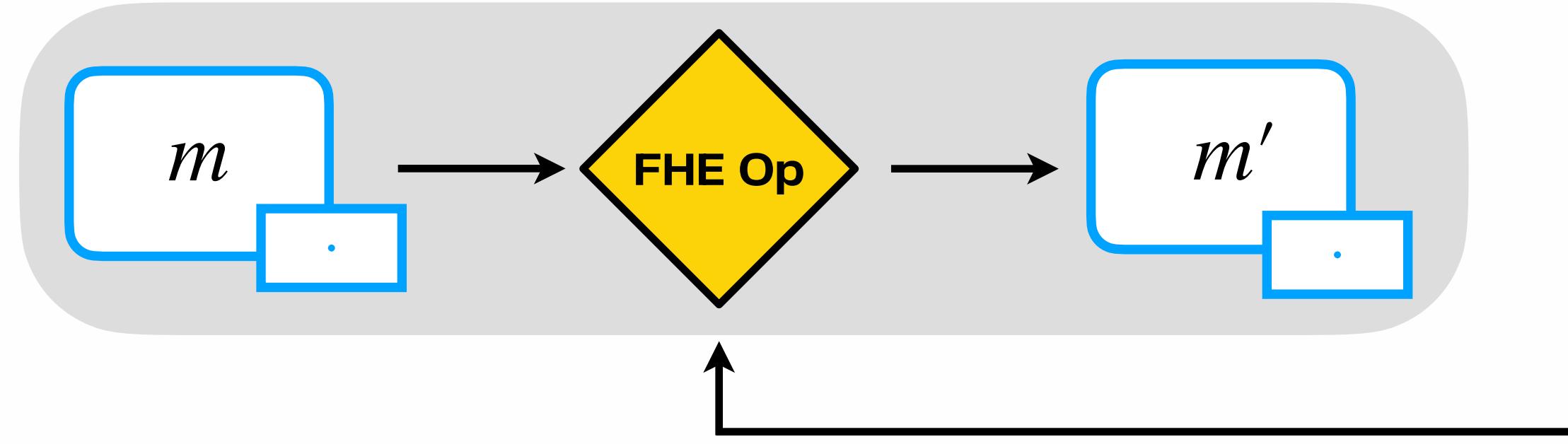
Shapiro-Francia* test

2.

Variance check

Check that $\text{Var}_{\text{exp}} \in \left[\frac{1}{2} \cdot \text{Var}_{\text{th}}, \text{Var}_{\text{th}} \right]$

Error Probability



Correct evaluation with probability $1 - P_{\text{fail}}$

$$\text{Ex: } P_{\text{fail}} = 2^{-40}$$

Rare event

On average, 2^{41} executions to observe **one** failure

Conditional Probability Basic Property

Importance Splitting*

$$\text{For } \tau_0 < \tau_1 < \dots < \tau_{J-1}, \Pr \left[|e| > \frac{\Delta}{2} \right] = \Pr \left[|e| > \frac{\Delta}{2} \mid |e| > \tau_{J-1} \right] \cdot \dots \cdot \Pr \left[|e| > \tau_1 \mid |e| > \tau_0 \right] \Pr \left[|e| > \tau_0 \right]$$

$\gg p_{\text{fail}}$

$\gg p_{\text{fail}}$

$\gg p_{\text{fail}}$

Parameter Oracle Bis

Client

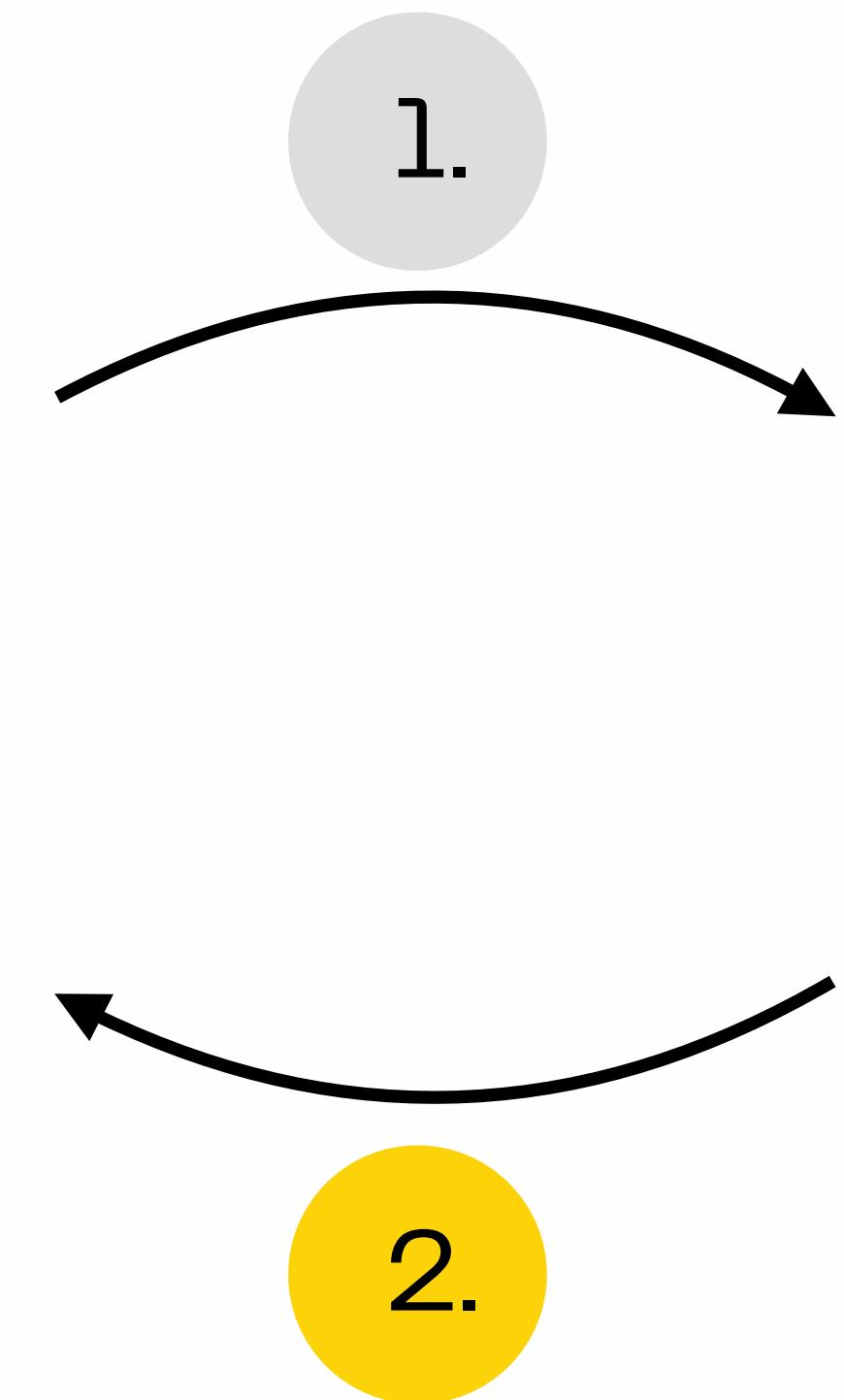
Message precision P

Linear operation number $\|\cdot\|_2$

Error Probability P_{fail}

Security Level λ

How to fix these last parameters?



Oracle

β_{KS}	ℓ_{KS}	β_{PBS}	ℓ_{PBS}
n	k	N	
σ_{LWE}	σ_{GLWE}		
χ_s	χ_e		
Δ	q		

Fixing the last parameters

Failure Probability P_{fail}

$$P_{\text{fail}} \leq 2^{-40} \iff \text{KS} \rightarrow \text{PBS is correct}$$

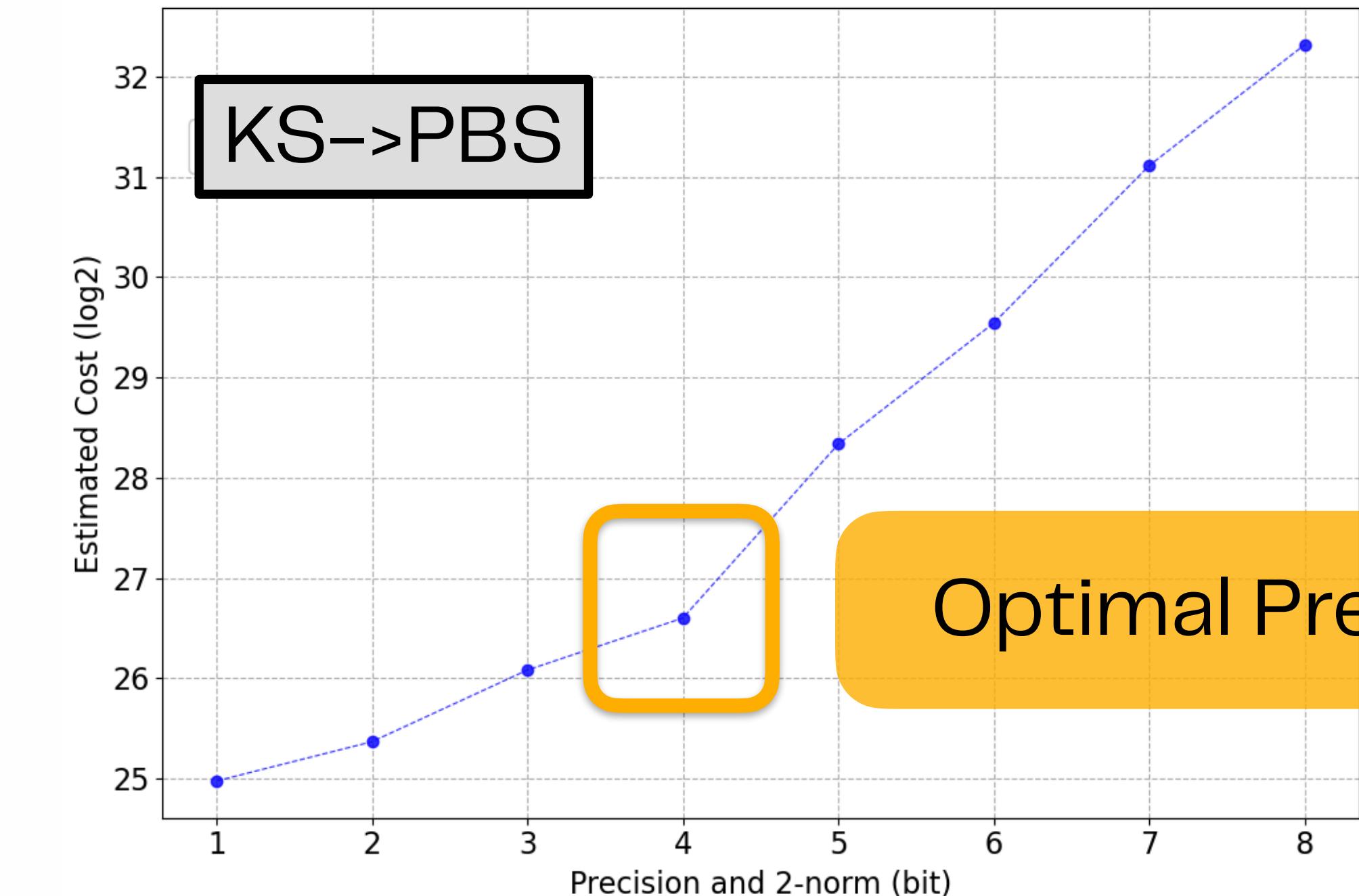
99.9999999991%

Linear Ops $\|\cdot\|_2$

$$\text{MsgMod, CarryMod} \Rightarrow \|\cdot\|_2 = \left\lfloor \frac{\text{MsgMod} \cdot \text{CarryMod} - 1}{\text{MsgMod} - 1} \right\rfloor$$

Security λ

Default value: $\lambda = 128$

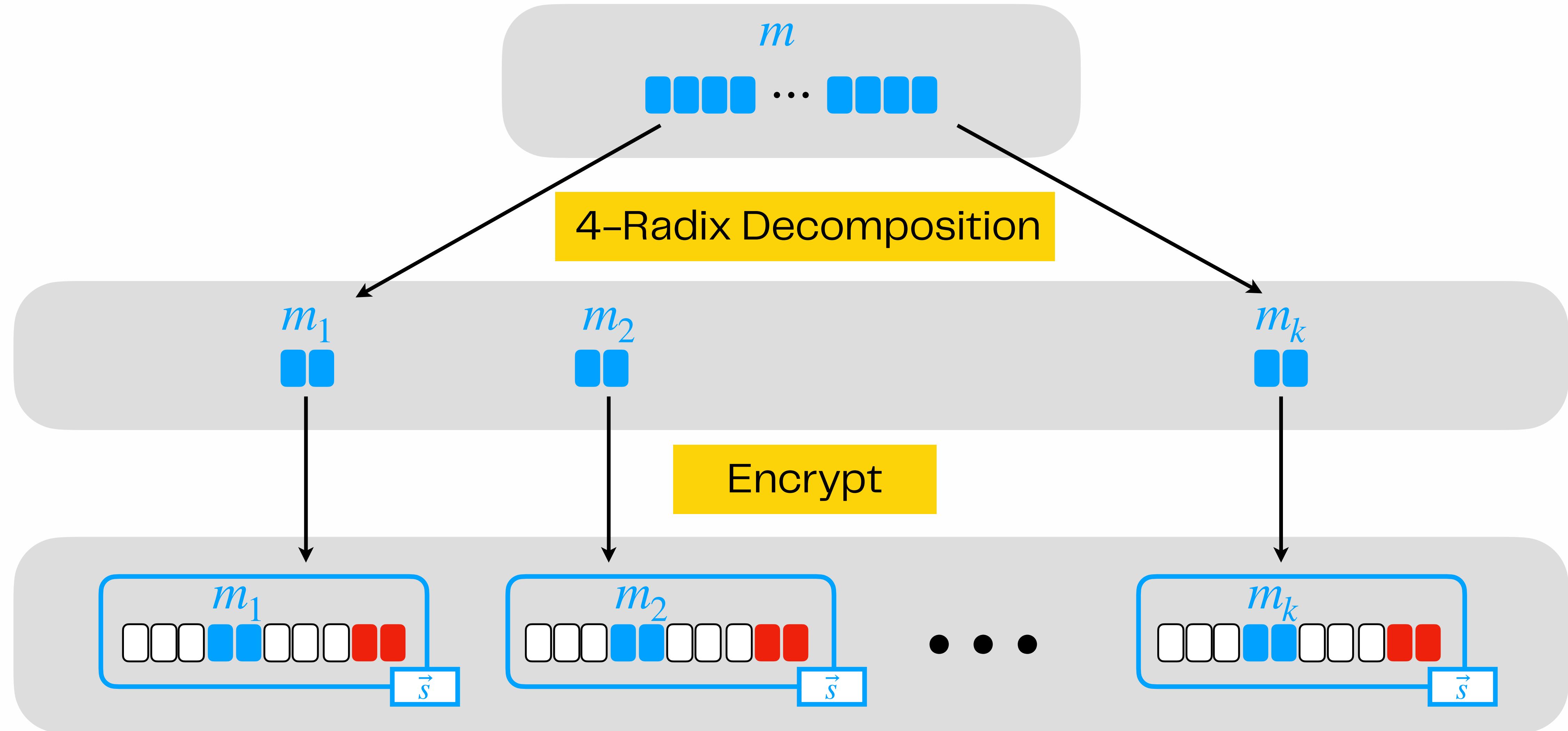


Precision P

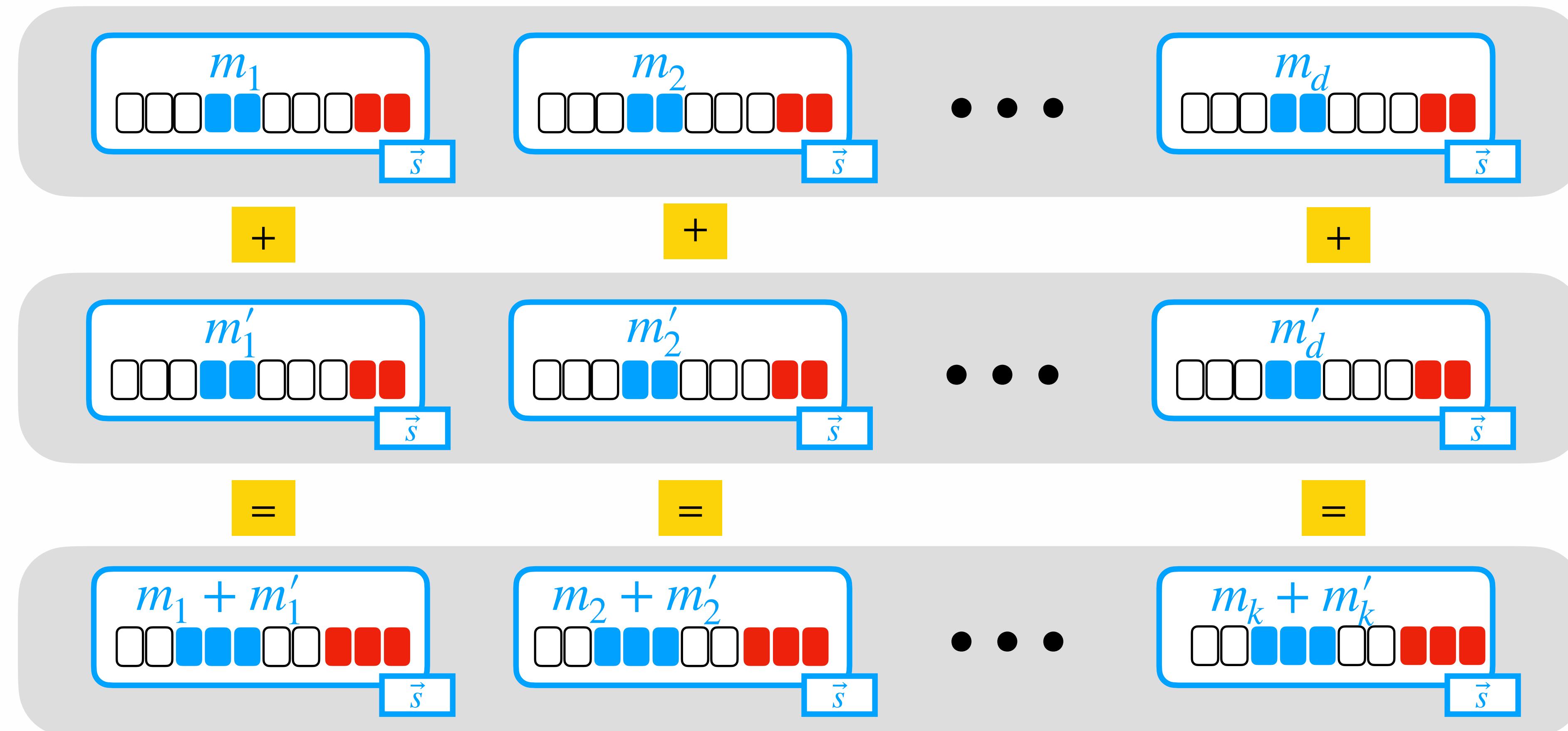
$\text{MsgMod} = 2^2$ & $\text{CarryMod} = 2^2$

Building Homomorphic Integers

Integers

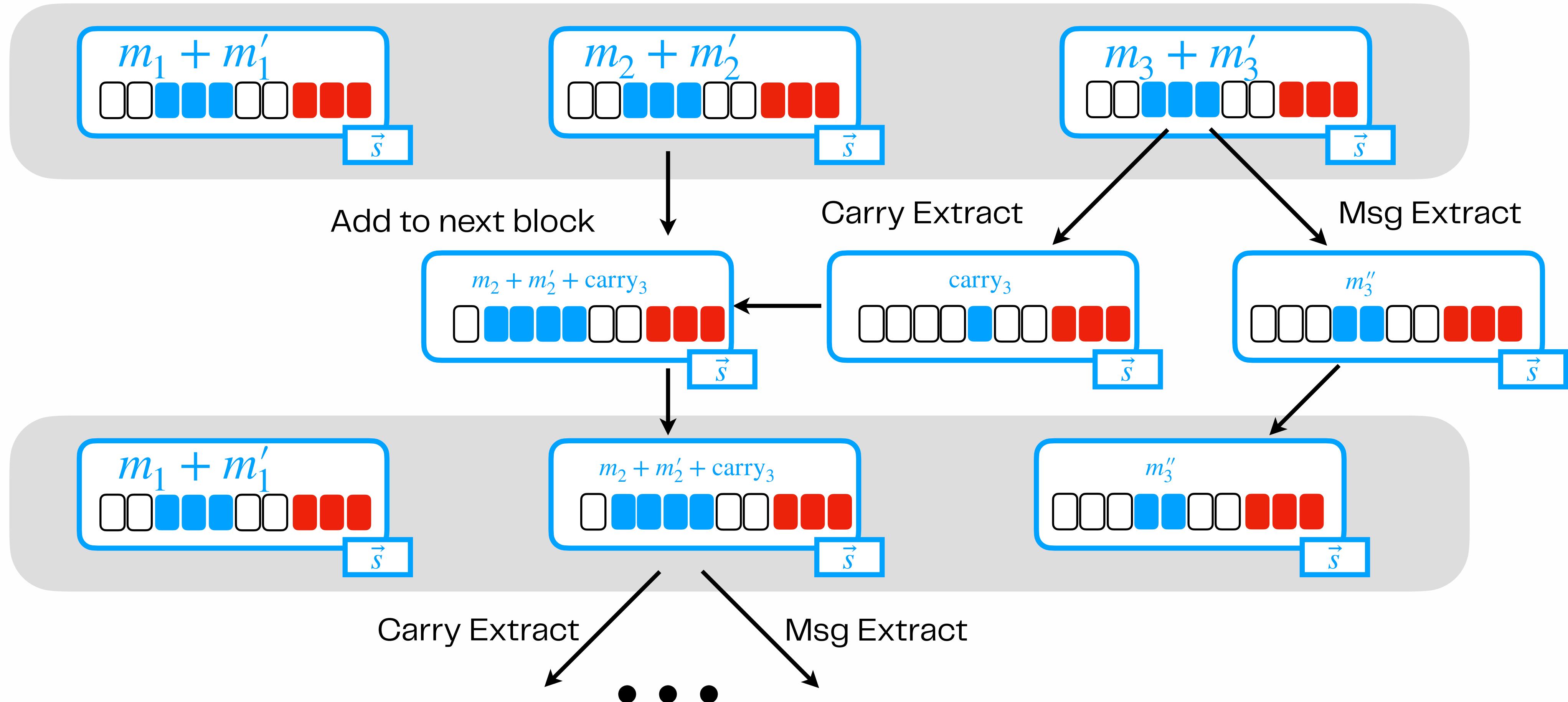


Addition



Output Encoding \neq Input Encoding \Rightarrow Needs to propagate carries

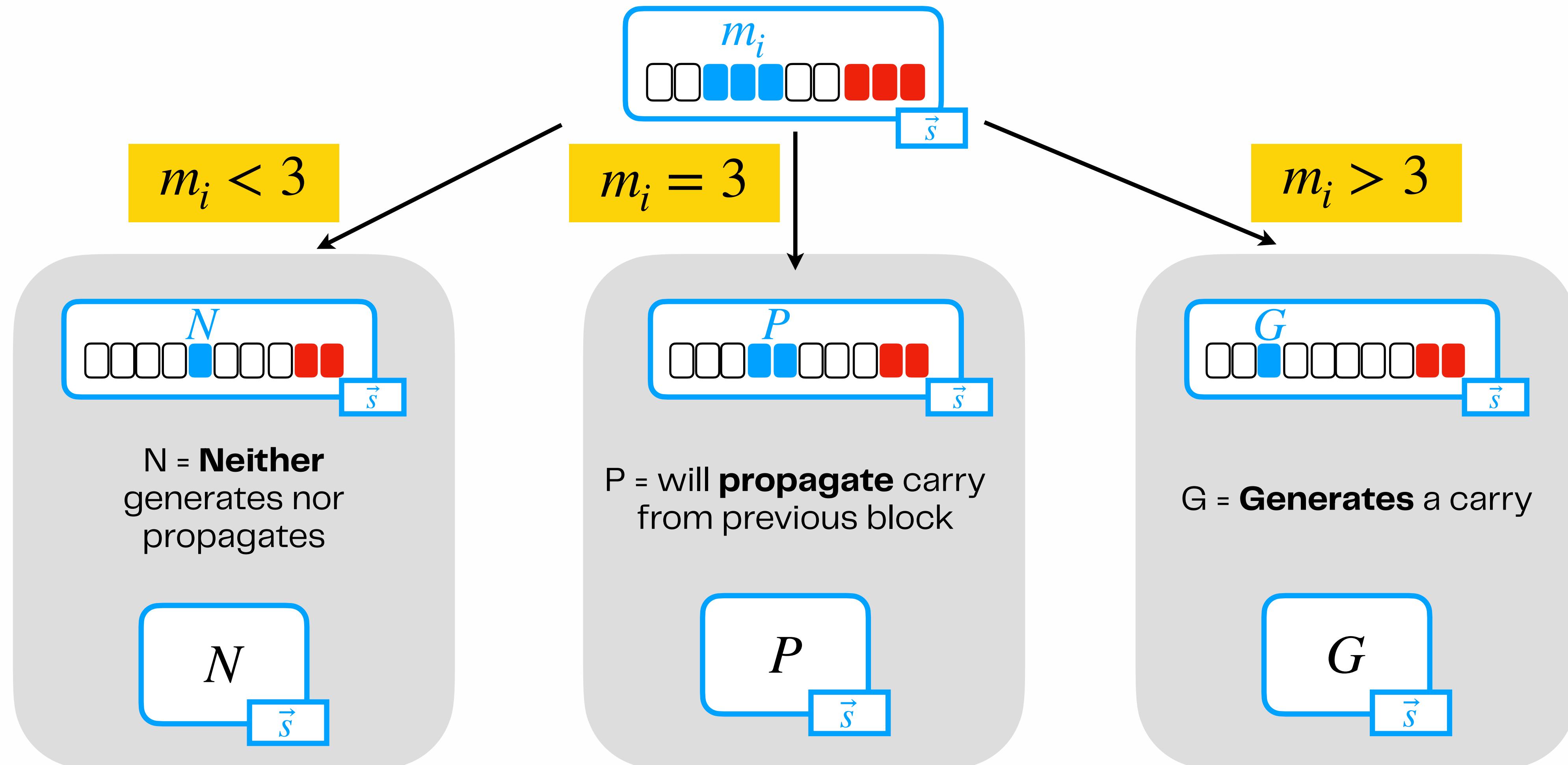
Propagating Carries



Highly sequential: $O(\text{BlockNumber})$

Faster Carry Propagation 1/2

1. Re-encode the state of each block



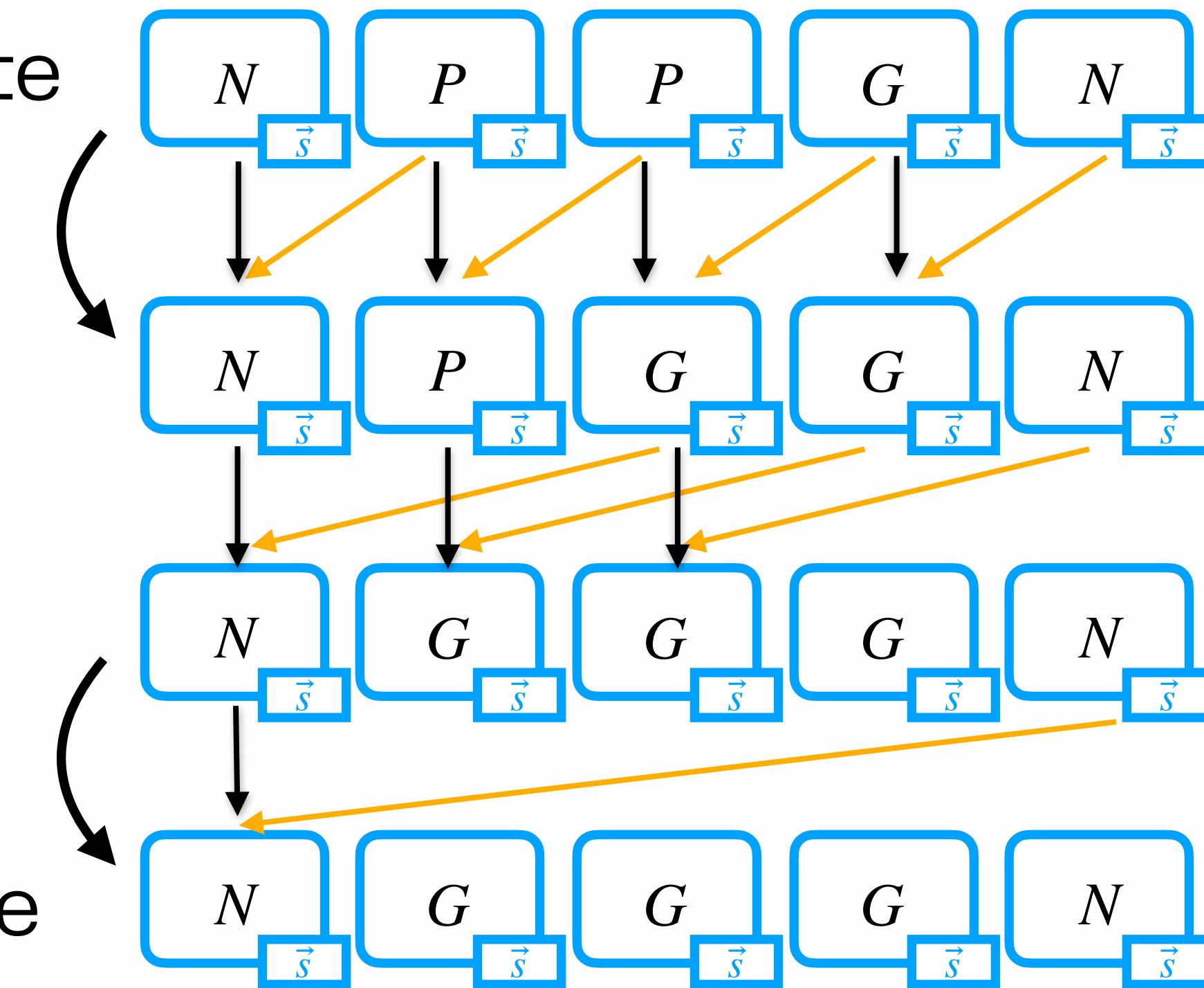
Faster Carry Propagation 1/2

2. Resolve the final state of the 'P' blocks [Hillis/Steele*]

		Output
		↓ ↗
N	N/P/G	N
G	N/P/G	G
P	N	N
P	P	P
P	G	G

Truth Table

Initial State



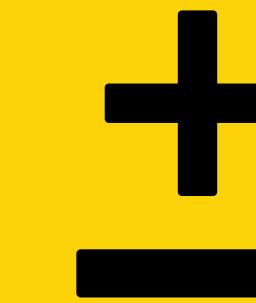
Final State

Leverage Parallelism: Depth $O(\log_2(\text{BlockNumber}))$

Other operations



Operators:
 $\&$ & $\|$, \ll , \gg , $=$, \leq , \geq , ...



Signed Integers
(2's complement)



Complete Arithmetic
+ - \times /



Operations with
overflow flag

Benchmarking FHE & Timings

How to benchmark FHE?

Security

$\lambda = 128$ bits

Error Probability

$P_{\text{fail}} = 2^{-40}$

A graph of FHE operators
s.t. Input Encoding == Output Encoding

KS → PBS with Integer encoding

MsgMod, CarryMod, $\| \cdot \|_2$

MsgMod = 2^2 , CarryMod = $2^2 \Rightarrow \| \cdot \|_2 = 5$

Hardware

AMD EPYC 9R14 CPU @ 2.60GHz

Latency/Throughput

Latency

Some Benchmarks

KS→PBS

with $P_{\text{fail}} = 2^{-40}$
using TFHE-rs

Precision	2 bits	4 bits
Timings	6.09 ms	12.72 ms

Sequential Timings

Precision	2 bits	4 bits
Timings	3.94 ms	6.01 ms

Parallelized Timings (grouping factor 3)

AND Gate

with $P_{\text{fail}} = 2^{-80}$

Library	TFHE-rs v0.5	TFHElib (main)	OpenFHE v1.1.1 w/ Hexl
Timings	7.62 ms	19 ms	21.8 ms

Sequential Timings

Integer Operations

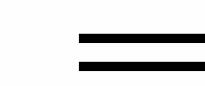
Operation	FheUint32	FheUint64
Add/Sub (+,-)	106 ms	81 PBS
Mul (x)	218 ms	481 PBS
Equal/Not Equal (eq, ne)	48.2 ms	19 PBS
Comparisons (ge, gt, le, lt)	83.8 ms	31 PBS
Max/Min (max, min)	120 ms	79 PBS
Bitwise operations (&, , ^)	17.7 ms	16 PBS
Div/Rem (/, %)	3.82 s	2761 PBS
Left/Right Shifts/Rotations (<<, >>)	135 ms	213 PBS

TFHE-rs timings with 192 cores

Conclusion

How can we abstract **FHE programming** from its cryptographic complexity to match the simplicity of **traditional coding**?

Trusted Automatic Parameter Generation
Efficient **Distribution** and **Failure** Tests



Foolproof Automatic
Parameter Selection

Static Parameters with the best
precision-complexity ratio



Efficient on a wide
range of use cases

Thank you.

ZAMA

Contact and Links

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zama.ai

[Github](#)

[Community links](#)

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